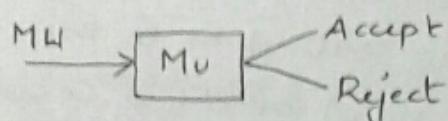


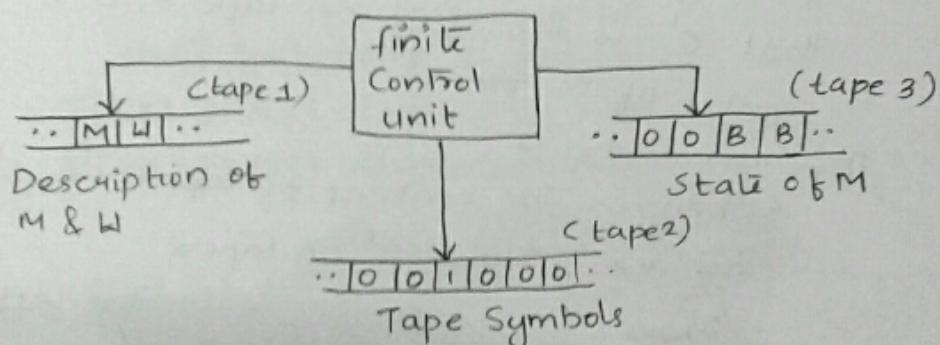
## Universal Turing Machine

- \* It is a TM for all other TM
- \* UTM is a multitape TM
- \* A reprogrammable TM called universal TM ( $M_u$ ) is an automaton that takes as input the description of any TM  $M$  and a string  $w$ .
- \*  $M_u$  can simulate the computation of  $M$  on  $w$ .



$$M_u = \{ \langle M, w \rangle / M \text{ is a TM and } M \text{ accepts } w \}$$

## Organisation of a Universal TM



UTM is a multitape TM consists of 3 tape.

- \* First tape holds description of  $M$  and input  $w$ .
- \* Second tape is used to hold the simulated tape of  $M$  using the same format has for the code of  $M$  ie, tape symbol  $x_i$  of  $M$  will be represented as 0's and tape symbols will be separated by '1'.
- \* Third tape of  $M_u$  hold the current state of  $M$ , in which the state  $Q_i$  is represented by  $i$  number of zeros.

The operations of Mu can be summarised as follows:

Step 1 :- Check code for M is valid for some TM 'M'. It not halt without accepting.

Step 2 :- Initialize second ~~steps~~ tape to contain the input w in its coded form

Eg:-

|            |   |    |     |
|------------|---|----|-----|
| 1/p        | 0 | 1  | B   |
| coded form | 0 | 00 | 000 |

Step 3 :- Place start state of M to the third tape. Then move the head of Mu's second tape to the first simulated cell.

Step 4 :- To simulate a move of M, Mu reaches on its first tape for a transition  $o^i | o^j | o^k | o^l | o^m$  such that  $o^i$  is the state on tape 3.  $o^j$  is the tape symbol of M.

This transition is

- changes the content of tape 3 to  $o^k$
- Replace  $o^j$  with  $o^l$  on tape 2
- Move head on tape 2 to the left/right depending on the value of  $o^m$ .

Step 5 :- If M has no transition that match simulated symbol and tape in step 4, M hold in the simulated configuration M.

Step 6 :- If M enters its accepting, then the Mu accept.